

UC/OSii 在 M251 上的移植

NuMicro® 32位系列微控制器范例代码介绍

文件信息

应用简述	在 M251 系列单片机上移植 UC/OSii_V2.93 的示例代码
BSP 版本	M251_M252_M254_M256_M258_Series_BSP_CMSIS_V3.02.003
开发平台	NuMaker-M258KE Ver1.1

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1. 概述

本代码介绍在 M251 系列单片机上移植 UC/OSii 需要添加哪些文件。

代码中的工程建立了相互通信的两个任务。

1.1 原理

为何要用 UC/OS? 因为用了 UC/OS,每个任务都像是独占了 CPU，因此，一个应用可以分成几个任务，而这些任务，可以由不同的人编写。

当一个任务被另一任务打断时，就像是 CPU 去执行中断代码了，过一会儿 CPU 就会回来。每一个任务都有独立的栈区，在任务开始调度前，每个栈区被初始化成像是刚刚发生中断一样，CPU 寄存器被压入了栈区。

```
S_STK *OSTaskStkInit(void (*task)(void *p_arg), void *p_arg, OS_STK *ptos, INT16U opt)
{
    OS_STK *stk;

    (void)opt; /* 'opt' is not used, prevent warning */
    stk = ptos; /* Load stack pointer */

    /* Registers stacked as if auto-saved on exception */
    *(stk) = (INT32U)0x01000000uL; /* xPSR */
    *(--stk) = (INT32U)task; /* Entry Point */
    *(--stk) = (INT32U)OS_TaskReturn; /* R14 (LR) */
    *(--stk) = (INT32U)0x12121212uL; /* R12 */
    *(--stk) = (INT32U)0x03030303uL; /* R3 */
    *(--stk) = (INT32U)0x02020202uL; /* R2 */
    *(--stk) = (INT32U)0x01010101uL; /* R1 */
    *(--stk) = (INT32U)p_arg; /* R0 : argument */

    /* Remaining registers saved on process stack */
    *(--stk) = (INT32U)0x07070707uL; /* R7 */
    *(--stk) = (INT32U)0x06060606uL; /* R6 */
    *(--stk) = (INT32U)0x05050505uL; /* R5 */
    *(--stk) = (INT32U)0x04040404uL; /* R4 */

    return (stk);
}
```

建好所有的任务并且栈区初始化好以后，CPU 开始执行最高优先级任务代码，就像是刚从中断返回来一样，这个工作是在函数 OSStart(), OSStartHighRdy() 和 PendSV_Handler 中做的。

任务的切换在函数 `PendSV_Handler()` 中。最高优先级任务栈区的首地址被读出来写入栈指针 `PSP`，然后执行中断返回，CPU 就返回到最高优先级任务去执行代码了。

```
PendSV_Handler
    CPSID    I                ; Prevent interruption during context switch
    MRS      R0, PSP          ; PSP is process stack pointer

    CMP      R0, #0
    BEQ      OS_CPU_PendSVHandler_nosave    ; equivalent code to CBZ from M3 arch to M0 arch
                                           ; Except that it does not change the condition code flags

    SUBS     R0, R0, #0x10    ; Adjust stack pointer to where memory needs to be stored to
                               ; avoid overwriting
    STM      R0!, {R4-R7}    ; Stores 4 4byte registers, default increments SP after each
                               ; storing
    SUBS     R0, R0, #0x10    ; STM does not automatically call back the SP to initial
                               ; location so we must do this manually
    LDR      R1, =OSTCBCur    ; OSTCBCur->OSTCBStkPtr = SP;
    LDR      R1, [R1]
    STR      R0, [R1]        ; R0 is SP of process being switched out

                               ; At this point, entire context of process has been saved
OS_CPU_PendSVHandler_nosave
    PUSH     {R14}           ; Save LR exc_return value
    LDR      R0, =OSTaskSwHook ; OSTaskSwHook();
    BLX      R0
    POP      {R0}
    MOV      R14, R0

    LDR      R0, =OSPrioCur   ; OSPrioCur = OSPrioHighRdy;
    LDR      R1, =OSPrioHighRdy
    LDRB     R2, [R1]
    STRB     R2, [R0]

    LDR      R0, =OSTCBCur     ; OSTCBCur = OSTCBHighRdy;
    LDR      R1, =OSTCBHighRdy
    LDR      R2, [R1]
    STR      R2, [R0]
    LDR      R0, [R2]          ; R0 is new process SP; SP = OSTCBHighRdy->OSTCBStkPtr;
    LDM      R0!, {R4-R7}     ; Restore R4-R7 from new process stack
    MSR      PSP, R0          ; Load PSP with new process SP
    MOV      R0, R14
    MOVS     R1, #0x04         ; Immediate move to register
    ORRS     R0, R1            ; Ensure exception return uses process stack
    MOV      R14, R0
    CPSIE    I
    BX       LR                ; Exception return will restore remaining context

    ALIGN
    END                        ; Ensures that ARM instructions start on four-byte boundary
```

PendSV 中断代码的流程图见 图 1-1. 进入 PendSV 中断后, 如果 PSP 不等 0, 表明 CPU 刚从一个任务跳转过来, 要先保存当前任务的上下文, 包括栈指针 SP, 然后再获取新任务的上下文。

PSP=0 表明刚从 OSStart()开始调度任务, 不需要保存当前任务的上下文。读出当前最高优先级任务的上下文, 栈指针 PSP 指向最高优先任务的栈区, 然后执行中断返回, CPU 就去执行最高优先级任务去了。

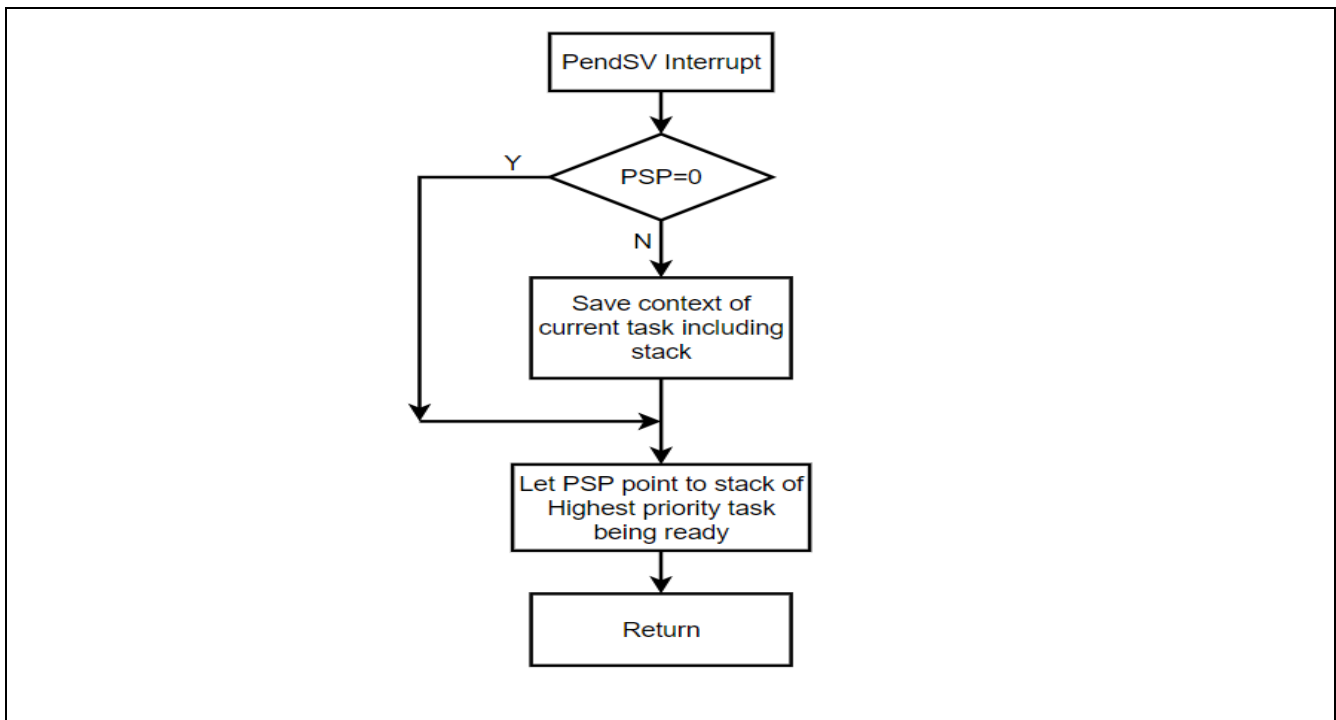


图 1-1 PendSV 中断代码流程图

1.2 执行结果

任务 2 让一个 LED 闪烁, 闪烁间隔或周期由任务 1 发过来。

2. 代码介绍

工程结构见图 2-1，工程中只包含必要文件。

由于版权问题，软件包中不含UC/OS内核代码文件，请到相关网站下载US/OS内核代码，按如下目录添加到工程中即可。

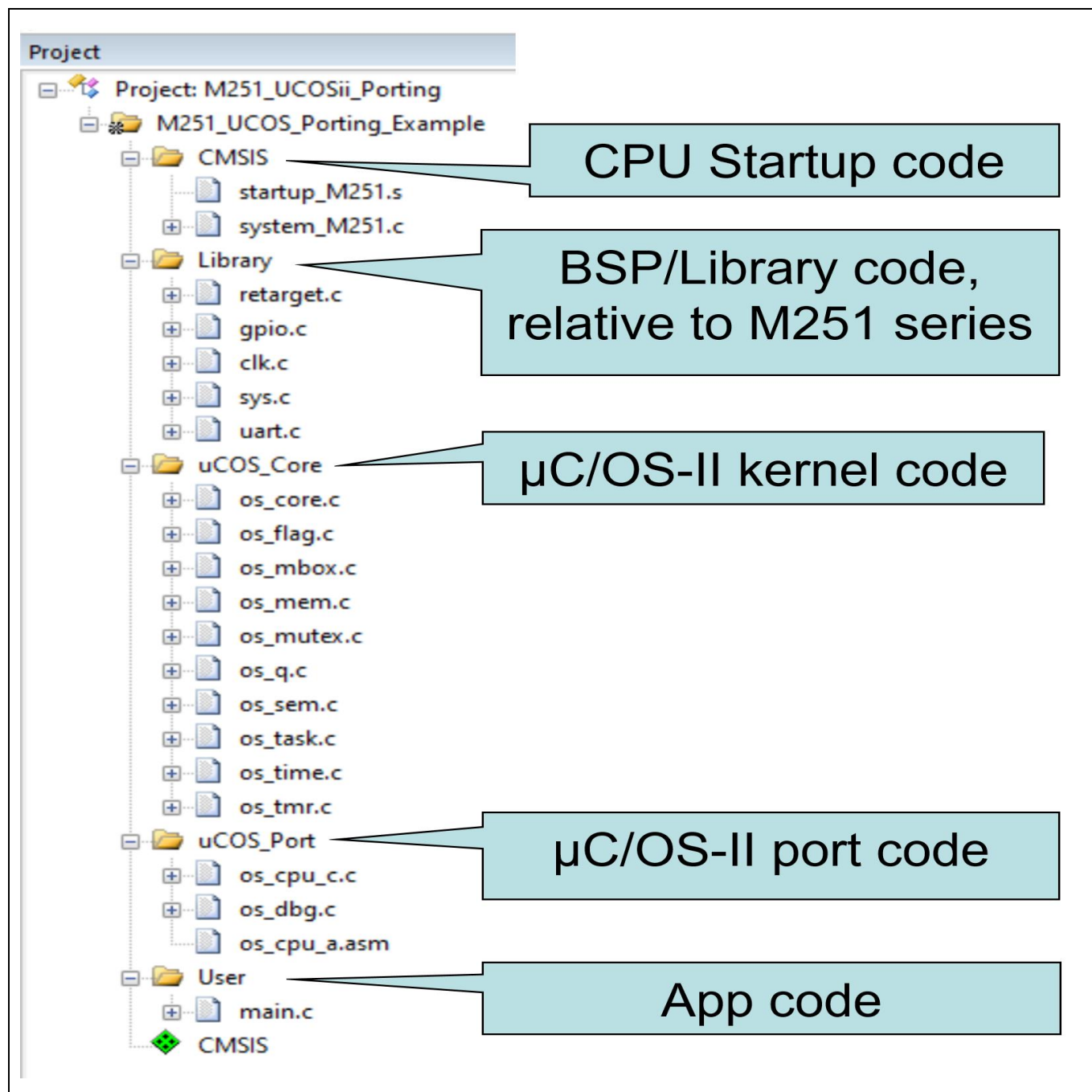


图 2-1 工程结构

M251 的各外设初始化后，先建立信号量、邮箱和队列，然后建立任务。执行 OSStart()之后，代码开始任务调度。

```
int main(void)
{
    int32_t Cnt32;

    SYS_UnlockReg();          /* Unlock protected registers */
    SYS_Init();               /* Init System, peripheral clock and multi-function I/O */
    SYS_LockReg();            /* Lock protected registers */

    UART0_Init();             /* Init UART0 for printf */
    printf("\n\nCPU @ %u Hz\n", SystemCoreClock);

    OSInit();                 // uCOS initialization.
    p_MailBox_1 = OSMboxCreate((void *)0); // create a mail box

    OSTaskCreateExt((void(*) (void *))task1_Highest, // create a task
                    (void *)0,
                    (OS_STK *)&task1_stk_H[LED_TASK_STK_SIZE - 1],
                    (INT8U)TASK1_PRIO_Highest,
                    (INT16U)TASK1_PRIO_Highest,
                    (OS_STK *)&task1_stk_H[0],
                    (INT32U)LED_TASK_STK_SIZE,
                    (void *)0,
                    (INT16U)OS_TASK_OPT_STK_CHK | OS_TASK_OPT_STK_CLR | OS_TASK_OPT_SAVE_FP);

    OSTaskCreate((void(*) (void *))task2_LED, // create another task
                 (void *)0,
                 (OS_STK *)&task2_stk_led[LED_TASK_STK_SIZE - 1],
                 (INT8U)TASK2_PRIO_LED);

    OSStart();                // Start schedule

    while (1) {} ;            // This sentence would be never executed.
}
```

每个任务都是一个 `while(1){ }` 循环。在循环中必须调用系统函数，如 `OSTimeDaly()`, `OSMboxPost()`等，如此，其它任务才有时间执行。

系统节拍定时器必须在最高优先级的任务里启动，否则，如果在任务调度之前发生了节拍定时器中断，系统可能崩溃。

任务 1 优先级最高，开头启动 `SysTick` 定时器。`while(1)`循环里，每 3 秒钟给任务 2 发一次数据。

```
void task1_Highest(void *pvParameters)
{
    SysTick_Config(SystemCoreClock / OS_TICKS_PER_SEC);    // MUST active tick-timer in highest task
                                                           // active tick-timer

    while (1)
    {
        OSTimeDlyHMSM(0, 0, 3, 0);                        // Delay 3 seconds

        if (Mail_Message[0] != 100) Mail_Message[0] = 100;
        else Mail_Message[0] = 500 ;                      // Configure message in mail box

        OSMboxPost(p_MailBox_1, Mail_Message);            // Send messgae via mail box
    }
}
```

任务 2 是让 LED 闪烁，闪烁周期，按任务 1 发送过来的时间。

```
void task2_LED(void *pvParameters)
{
    int32_t *ptr ;
    int32_t static Time = 300 ;

    while (1)
    {
        ptr =OSMboxAccept(p_MailBox_1); //Check whether there are messages in the mailbox

        if (ptr !=(void*)0) Time = *ptr; //Read out if there is any message in the mailbox

        OSTimeDlyHMSM(0, 0, 0, Time);    // Delay time according to message.
        PB14 ^= 1 ;                      // LED flash
    }
}
```

3. 软件与硬件需求

3.1 软件需求

- BSP 版本
 - M251_M252_M254_M256_M258_Series_BSP_CMSIS_V3.02.003
- IDE 版本
 - Keil uVersion 5.26

3.2 硬件需求

- 电路组件
 - NuMaker- M258KE Ver1.1

4. 目录信息









 EC_ M251_UCOSii_Porting_V1.00	
 Library	Sample code header and source files
 CMSIS	Cortex [®] Microcontroller Software Interface Standard (CMSIS) by Arm [®] Corp.
 Device	CMSIS compliant device header file
 StdDriver	All peripheral driver header and source files
 SampleCode	
 ExampleCode	Source file of example code
 uCOS_II	Files relative to uCOS

图 4-1 目录信息

5. 范例程序执行

1. 根据目录信息章节进入 ExampleCode 路径中的 KEIL 文件夹，双击 *M251_UCOSii_Porting.uvprojx*。
2. 进入编译模式界面
 - 编译
 - 下载代码至内存
 - 进入 / 离开仿真模式
3. 进入仿真模式界面
 - 执行代码

6. 修订纪录

Date	Revision	Description
2022.07.27	1.00	初始发布。

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